

# The ATA Commensal Observing System

## ATA Memo #89

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### Abstract

This memo describes the system used to conduct commensal correlator and beamformer observations at the ATA. This system was deployed for  $\sim 2$  years until the ATA hibernation in 2011 and was responsible for collecting  $>5$  TB of data during thousands of hours of observations. The general system design is presented (§2) and the implementation is discussed in detail (§3). I emphasize the rationale for various design decisions and attempt to document a few aspects of ATA operations that might not be obvious to non-insiders. I close with some recommendations (§4) from my experience developing the software infrastructure and managing the correlator observations. These include: reuse existing systems; solve, don't avoid, tensions between projects, and share infrastructure; plan to make standalone observations to complement the commensal ones; and be considerate of observatory staff when deploying new and unusual observing modes. The structure of the software codebase is documented (§3, §A, §B).

## 1 Introduction

One of the major design goals of the Allen Telescope Array (ATA; [Welch et al., 2009](#)) was the ability to share observatory resources to conduct multiple independent observing programs simultaneously — a feat generally referred to as commensal observing.<sup>1</sup> This memo describes the system used

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<sup>1</sup>In biology, commensalism is a symbiotic arrangement between mutualism and parasitism: it is the coexistence of two organisms in which one benefits and the other neither

to conduct the largest commensal observing campaigns before the ATA’s hibernation in 2011. The primary campaign was a joint survey of regions in the Galactic plane: a traditional radio-astronomical survey looking for broadband Galactic radio transients (*à la* Hyman et al., 2005) and a SETI (Search for Extraterrestrial Intelligence) search for narrowband sources. In both cases, the Galactic population is clearly best probed by looking in the plane rather than searching the sky uniformly. Various components of the traditional search have been referred to as the GCS (Galactic Center Survey), the AGCTS (ATA Galactic Center Transient Survey), gal90 (Galactic plane,  $l = 90^\circ$  region), or the Kepler or Cygnus X-3 surveys. These are now all grouped together as AGILITE, the ATA Galactic Lightcurve and Transient Experiment. (The survey description paper is currently in preparation.) Because of the low overhead to doing so, commensal correlator observations were also made during the SETI “exoplanets” and “waterhole” surveys, but there are currently no specific plans to use the correlator data from these undertakings.

The original vision that motivated the goal of substantial commensal observing on the ATA was one in which scheduling was based on traditional radio astronomy applications but SETI searches ran continuously in the background as well. The traditional observing program would drive the telescope, controlling the pointing and most hardware systems, and use the ATA’s correlator backends to take data. Meanwhile SETI observations would be performed using the ATA’s beamformer backends more-or-less passively, choosing any promising observation targets lying within whatever field of view (FOV) was chosen by the traditional observing program. It’s worth noting that this is not the only commensal observing scheme that might feasibly be implemented at the ATA. For instance, with a sufficiently large number of dishes and two relatively undemanding traditional radio astronomical observing projects, one could partition the array into two subarrays and run the two projects simultaneously, each using separate antennas and correlators. (One could argue that this is in fact not a commensal mode since so few resources would be shared between the two projects.)

The earliest commensal observations at the ATA were performed in August 2008 by Karto Keating and Tom Kilsdonk, but these and a few other efforts never became routine observing modes. As such, I refer to the sys-

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benefits nor suffers. It derives from the Latin *cum mensa*, “sharing a table,” originally alluding to the sharing of food scraps in particular.

tem under discussion as “the ATA commensal observing system,” without further qualification. A design for an observatory-wide commensal observing system, and a deeper discussion of the observatory-wide software systems, is presented in [Gutierrez-Kraybill et al. \(2010\)](#).

## 2 Survey and System Design

Observations for the commensal campaigns are scheduled in blocks as per the standard system in use at the ATA. In contrast to the long-term vision of ATA commensal observations, it is SETI, not the traditional radio observing program, that’s “in the driver’s seat” for the observations: SETI software takes responsibility for all of the telescope hardware, most importantly the pointing. This arrangement came about because SETI already had a well-established observing system called Prelude, which had been adapted from its Arecibo roots to work at the ATA as well. Given this existing codebase, the project was approached with the plan of minimizing the amount of changes required to Prelude, while adding a separate “commensal observer” component that would take care of everything related to the commensal correlator campaigns.

With Prelude in charge of pointing the antennas and taking care of SETI’s data-taking, the responsibilities of the commensal component of the observing campaigns are *extremely* constrained: essentially, all it can do, and all it needs to do, is turn the correlators on and off. One could envision a much more complex system in which the commensal observer dynamically notifies Prelude of various needs (“please visit a phase calibrator”), but the system is vastly simplified if all such decision-making is centralized in Prelude. This simplification was made possible by pre-negotiating such decisions as pointing centers, dwell times, calibrator intervals, and focus settings.

In principle, the commensal observer could be completely ignorant of Prelude and its workings; by monitoring the current pointing directions of the ATA dishes, it could decide when a source was being observed and when the array was slewing, and in the former case it could do the necessary work to take data. It turns out, however, that “where is the ATA pointing?” and “are we tracking or slewing?” are questions that are more difficult to answer than one might think: obtaining ATA antenna telemetry requires fairly complex code (cf. the implementation of the `atastatus` command), and glitches in the data stream make it difficult to interpret the data robustly.

Prelude knows exactly what its intentions are, however, so it was decided that the commensal observer would monitor a telemetry stream from Prelude to obtain pointing information. This stream is described more fully below (§3.3).

An important thread running through the design of the commensal system is automation: with thousands of hours of commensal observations scheduled, it's desirable to execute them using as little human intervention as possible. The existing Prelude system fortunately dealt with the difficult task of choosing targets and planning a night's observing.

### 3 System Implementation

The ATA commensal observing system is composed of a group of software tools and practices for using them. As mentioned above, some aspects of the commensal system were implemented in Prelude and its successor, SonATA (“SETI on the ATA”). This code is internal to SETI and is not discussed here.

Although the rest of the commensal software runs on a diverse set of hosts in a diverse range of situations (*e.g.*, at HCRO during observing; at UCB during data reduction), the whole codebase is maintained in a single Git (<http://git-scm.com/>) repository by the author. This document describes Git commit 88263be44c3d724e700e46b156e9b1dcfd0b1089, made on 2011 November 17. (Due to the nature of the Git version control system, this one value identifies an exact snapshot of the commensal source tree as well as its entire revision history.) A public version of the commensal repository is currently available at <http://astro.berkeley.edu/~pkwill/repos/commensal.git>. This URL will likely go stale as I am (hopefully...) soon leaving Berkeley. A public copy of the repository may be established at his GitHub account, <https://github.com/pkgw/>, and at a minimum the repository will be available upon request. (Once again, thanks to the design of Git, each repository stands alone and contains every version of every file as well as the complete revision history of the source tree.)

The vast majority of the source code is written in Python, with some tools written in shell script. The majority of the latter are Bourne shell scripts (most likely only compatible with `bash`), but a few are `tcs`h scripts, since the latter was the language used for most ATA observatory infrastructure. The repository also contains scheduling and analysis metadata.

The Python scripts reference some support modules for controlling the ATA that are distributed in the `mmm`<sup>2</sup> repository at <https://svn.hcero.org/mmm/pwilliams/pyata/>. A few secondary packages use NumPy, `miriad-python` (<http://purl.oclc.org/net/pkgwpub/miriad-python><sup>3</sup>) and/or my plotting package, `Omegaplot` (<https://github.com/pkgw/omegaplot/>). The system design is shown in schematic form in Figure 1.

### 3.1 Scheduling

As mentioned above, commensal observations are allocated blocks in the ATA schedule as per standard practice.

At the time of the observations, the observatory-wide system for enacting this schedule was fairly *ad hoc*. The main HCRO computing environment is a distributed Unix-type system, with user logins and home directories shared over the network to a variety of workstations and servers in the signal processing room (SPR). One user account, `obs`, has a shared password and by convention is used to conduct the vast majority of observatory operations. SETI software systems, however, are segregated from this shared system, and generally run on separate machines with distinct logins. I’m not familiar with their configuration.

Correlator observations are scheduled by constructing a large shell script in the `obs` account, `obs.run`, that serially invokes other shell scripts to perform standard observing tasks. Before the 2011 hibernation, most observatory time was devoted to correlator observations so this could be considered the standard ATA scheduling system. For purely SETI observations, the `obs.run` script idles for a fixed amount of time while the array is controlled from SETI computers.

Because the `obs.run` system proved to be unreliable and the commensal campaigns involved many identical observations, the commensal observer was designed to be scheduled and launched separately in a more automatic

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<sup>2</sup>The name derives from our weekly data analysis meeting, “MIRIAD Masterclass with Melvyn” (Wright).

<sup>3</sup>Note the unusual host name in this URL. This link is a “Permanent URL”, one that is intended to be stable over decade timescales and thus (hopefully) more appropriate for inclusion in the academic literature. This is accomplished merely by forwarding requests for the PURL to a changeable destination, somewhat like the URL shorteners that are currently popular. I’m not using PURLs for most of the links in this paper, but this one in particular already existed because of its publication in Williams & Bower (2010). See <http://purl.org/> for more information.

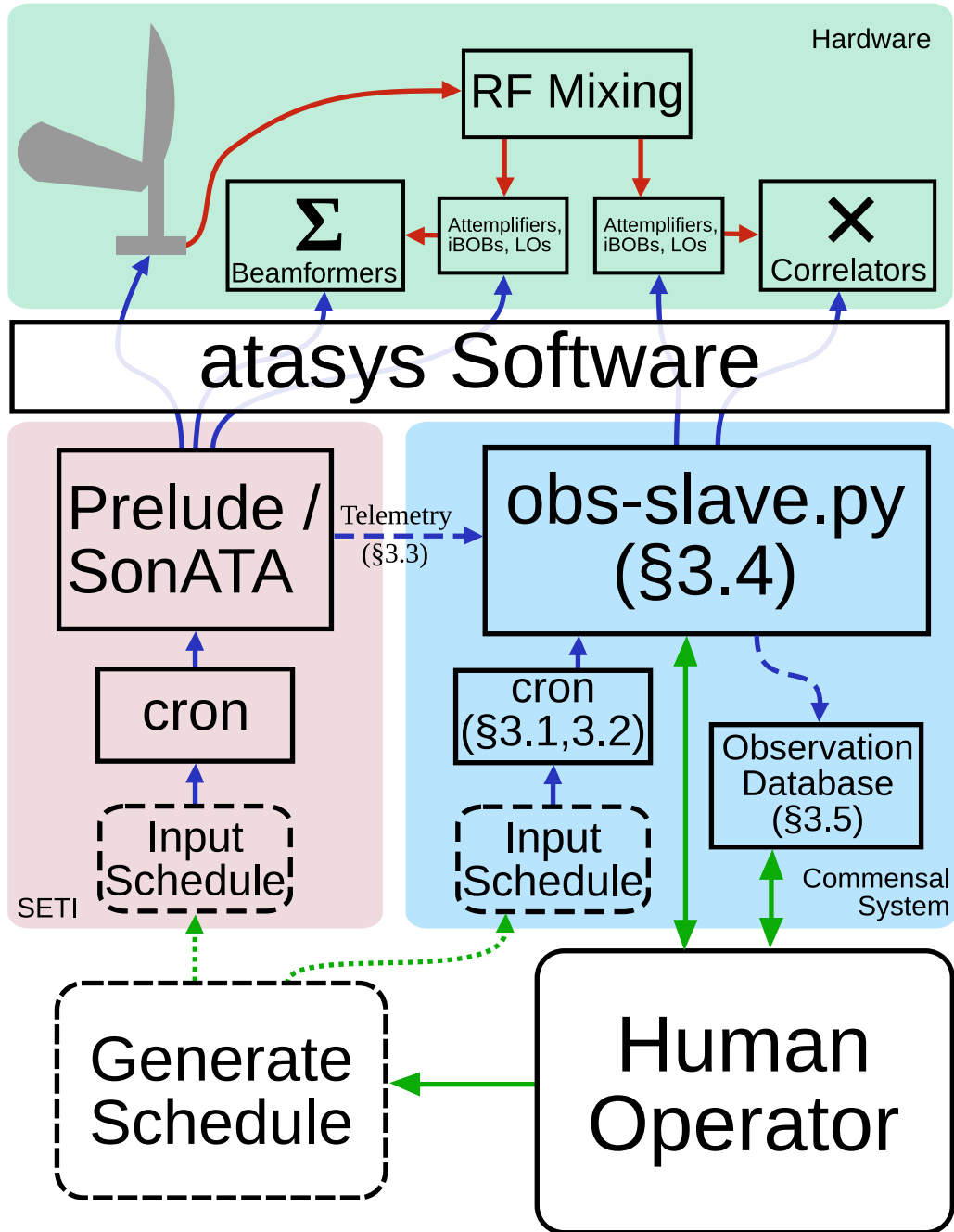


Figure 1: Schematic of the commensal observing system. See §3.

fashion. Commensal observing blocks are entered into a separate schedule file in the repository (`sched/current.tab`) using the commands `plan-schedsess` or, later, `plan-importtxt`. The observing processes are launched via `cron` (see below).

Meanwhile, the SETI observations are launched via a `cron` system as well in the SETI part of the network. This means that there are in effect at least *four* relevant versions of the ATA schedule: the textual/graphical version distributed to users, that schedule as incarnated in `obs.run`, the schedule used by the commensal observer, and the schedule used by Prelude. The large PiGSS (Pi GHz Sky Survey) transient survey also used its own scheduling system. The suboptimality of this situation should be obvious. Regarding the particular case of the commensal campaigns, however, things weren't too bad. While it was sometimes necessary to coordinate schedule changes with SETI staff over email, and mistakes could result in lost observing opportunities, most observations went smoothly. Rarely, miscommunications would result in `obs.run` and the commensal launcher attempting to observe simultaneously; due to poor lockouts in the observatory infrastructure, it's actually possible for both systems to operate at the same time, with expectedly poor results.

Each observing session scheduled in the commensal system was assigned a random 32-bit hexadecimal identifier, *e.g.* `8592a0ce`. This identifier was used for tracking data and metadata throughout the rest of the observing and analysis chain. (The original scheme, used in the 2009 GCS survey, used universally unique identifiers [UUIDs], which are a particular kind of 128-bit identifier. They are 36 characters long and were judged to be overkill.) Note that identifying a session by the day on which it occurs, for instance, is insufficient if there are multiple sessions in one day. In theory, it doesn't even work to identify a session by its start time, since the array could be partitioned into two subarrays with different sessions being executed during the same block of time.

### 3.2 Launching and Managing Observations

The main challenges in automatically launching commensal observations are robustness and control. The loose nature of the ATA online system makes it possible for the commensal observations to start up with the array in odd states, and there are inevitably network outages, computer crashes, hardware problems, *etc.* Meanwhile, the array schedule isn't written in stone, and the operator needs to be able to see what's going on and potentially alter array

operations at will.

Commensal correlator observations are launched automatically under the author's user account, `pkwill`. A launcher program, `obs-hourly`, runs hourly via `cron` (see `resources/obs.crontab`) on the host `obscntl` and consults the schedule file in a checkout of the repository found in `pkwill`'s home directory. If an observation is upcoming in the next hour, the launcher stays resident and `execs` the appropriate observing script at the right time. Because most observations are scheduled for at least a few hours, the hourly invocation of the launcher was intended to provide some insurance if the observing script crashes; this feature has been helpful only a handful of times.

Observations can also be “kickstarted” manually if the automatic invocation fails. This usually happens when an unusual hardware condition needs to be resolved before correlator observations can start. This is done with the `obs-kickstart` command, which uses GNU `screen` to launch the observing script in a detached, persistent virtual terminal. Both launchers have interlocks to attempt to prevent multiple observations from running simultaneously and to prevent observations from being launched outside of a scheduled commensal observing block.

Due to the distributed nature of the HCRO computing environment, observations can in theory be run from a variety of hosts. This capability was not usually exploited in ATA operations. Commensal observations were designed to be run from the `cron` host `obscntl` (AKA `auxcntl`) or, in early days, `tumulus`.

The `obs` user can monitor and control commensal observations with the command `commensal-mc`, which was set up to be easily accessible to the `obs` user and was intended to be straightforward to use. The command can perform diagnostics such as checking whether a commensal observing session is currently active (according to the commensal schedule), whether any commensal observing processes are currently active, killing or kickstarting observer processes, *etc.* Communication between processes and users is accomplished with special files in the shared ATA archive disk, `/ataarchive`, which is accessible to all HCRO computers and is writeable by most user accounts (notably, `obs` and `pkwill`). Commensal observing processes should be run as the user `pkwill`, which presents a difficulty if the user `obs` is attempting to kickstart a new process. This is dealt with by doing the kickstarting through a special passwordless SSH key with a preset command.

A more detailed look at the data coming out of a commensal observation can be obtained with `misc-parseplog`, which parses log output from the



Name	Code	Description
idle	0	Observing is not active, and the coordinates are undefined
slewing	1	The telescopes are slewing, and the coordinates specify their destination
pointed	2	The telescopes are pointed up on the specified coordinates

Table 1: Array states enumerated in the synchronization protocol.

commensal observer and does some simple diagnostics on the datasets being written by the data catchers. This program has proven to be quite helpful for checking target selection and catching occasional correlator failures which were both serious and undetected by the usual diagnostics.

### 3.3 Synchronization Protocol

As alluded to above, the commensal observing software monitors a telemetry stream from Prelude. This stream was designed to convey the minimal amount of information necessary for the commensal observer to easily drive the correlators. The stream is broken into messages, each of which transmits a timestamp, an array state, and observing coordinates. The states are listed in Table 1.

This architecture lends itself to a state-machine implementation in the commensal observer. The lack of state in the telemetry stream means that the observer and/or Prelude can die at arbitrary times and the observer will recover sensibly.

The telemetry stream is implemented with broadcast UDP packets on port 24243. (This port was chosen arbitrarily.) Each packet consists of 32 bytes: a 64-bit double-precision floating point Unix timestamp, a 32-bit unsigned integer state code, an unused 32-bit padding field, a double-precision right ascension coordinate measured in hours, and a double-precision declination coordinate measured in degrees. The values are big-endian and can be decoded with the Python `struct.unpack` specifier `>dIId`. The numerical codes corresponding to the states are also listed in Table 1. The specified transmission rate is one packet per second, though there’s nothing special setting this rate. Protocol decoding is implemented in `pylib/protocol.py`.

Clearly, it is not difficult to implement this protocol. The program `obs-`

`listener.py` can monitor the telemetry stream and print out low-level packet contents and check for unexpected state transitions. For example, it was used to diagnose a situation in which two Prelude processes were simultaneously running, one broadcasting a series of idle packets, the other attempting to run observations. The apparent oscillation between the idle and active states unsurprisingly confused the observer.

The program `obs-master.py` can be used to perform commensal-style observations when Prelude is inactive. Like Prelude, it performs target selection, drives the telescopes, and emits telemetry packets. Thanks to the clear division of responsibilities and the simplicity of the commensal observing protocol, the commensal observer can operate completely normally even though the software “in the driver’s seat” is completely different. This was also the experience when Prelude was replaced with SonATA, SETI’s next-generation observing system. Besides a change in the active port number (to allow simultaneous testing of Prelude and SonATA), the replacement of several racks of computer equipment and the entire software codebase was completely transparent to the commensal observer.

Compared to most elements of the commensal software package, the `obs-master.py` program is relatively complicated. It can either loop through a list of targets in order, with calibrator scans interspersed, or it can always observe the lowest-declination visible source. A few other parameters, such as the antenna focus position and interval between calibrator scans, are also configurable. It was not used extensively so its target-selection algorithms are relatively unsophisticated.

The current code to monitor broadcast UDP packets doesn’t work in the case that the emitter and listener are on the same host, which corresponds to situations when `obs-master.py` is being used. I suspect that this could be fixed with a better understanding of the proper low-level socket setup, but since this problem arises during manually-arranged observations, during those times I just tweak the socket code to skip actual network communication (by using the loopback device). This “solution” did lead to problems once or twice when I forgot to undo the tweak once observations were over.

### 3.4 Correlator Observer

The commensal correlator observer, `obs-slave.py`, is the most complex element of the commensal observing system. It is a single program that is launched by the hourly cronjob, monitors the telemetry stream, and operates the ATA

correlators to take data for transient searches.

On a low level, the commensal observer is responsible for driving a specific set of subsystems:

- the ATA correlators
- the amplifiers used by those correlators
- the local oscillators (LOs) used by those correlators
- the fringe-rotator iBOBs

The observer must take appropriate action based on state transitions in the telemetry stream or other external inputs (*e.g.*, abort commands).

The observer must also pay attention to some secondary systems to support its observations:

- It must have access to a catalog (mapping from coordinates to source names) so that correct target names can be shown to the operator and embedded in output datasets.
- It must generate ATA “ephemeris files” to be used by the fringe rotators and correlator data-catchers.
- It must check for abort signals, which is accomplished by monitoring the `/ataarchive` filesystem for the appearance of special files.
- It must generate detailed logs, since one wants to be able to debug subtle problems, check that hardware is being driven correctly, search for efficiency losses, and monitor the array.
- It must save appropriate state so that if it crashes and is restarted, the array hardware is reset *or not reset* so that datasets from before and after the crash may be calibrated consistently.

In order to accomplish all this, the observer is a multithreaded program with a main coordinator thread and various subthreads responsible for the subsystems. Because it is important that the observer be resilient to crashes, there’s also complex code to deal with Python exceptions as robustly as possible. The final program is still less than 700 SLOC (statement lines of code), a nice demonstration of the concision of Python. Looking at the source will confirm, however, that the code is quite dense.

The commensal observer stores each individual scan in a separate MIRIAD dataset simply named with a serial number. This approach makes a directory listing relatively opaque, but was hoped to be more robust than the standard ATA practice of naming datasets with some combination of source name, observing frequency, *etc.*, and appending to these datasets when multiple scans occur. Most MIRIAD tasks can stream through multiple datasets, but none can truly analyze only a portion of a dataset, so it should be more efficient to create many datasets and sometimes deal with them in bulk, rather than to create large datasets and sometimes subdivide them. There have also been instances where a large dataset has been rendered invalid due to some problem at the very end of an observation, and subdividing datasets helps minimize the damage incurred in these cases.

The datasets generated by the commensal observer are augmented with a few extra data items. A file within the dataset named `c_instr` records the identifier of the ATA correlator used to generate the data; there is perhaps an oversight in the ATA data catcher software that this information is not recorded in datasets otherwise. (This information is useful because failures in the digital electronics have correlator-specific effects.) Another file named `c_tstop` records the timestamp at which Prelude reported leaving the current pointing for another target — there will be a delay between Prelude issuing the slew command and the correlator data catchers shutting down and closing their datasets, so there may be some bad data at the end of a scan taken while the antennas are slewing.

One particular challenge faced by the commensal observer is that the standard ATA data catcher program, `atafx`, was designed to be run with a predefined integration time. By the nature of the commensal observing system, however, the commensal observer does not know how long each integration will last. (If the expected integration time were transmitted from Prelude to the observer, one would still have to check that it was honored, and dealing with unexpected circumstances would require all of the flexibility needed by the current implementation.) Given the current software, the best solution is actually to SSH to `strato`, the data-taking host, and kill the `atafx` processes. A project was started to change the data takers to be always-running servers so that integration could be stopped and started quickly and on-demand, but that code never reached deployment.

### 3.5 Observation Database

The commensal observing system includes not only software for making observations but also a set of tools for understanding the observations that have been made. These are built on a database of observing sessions and individual scans.

These post-observation tools could plausibly have been separated into a different software package, and that might arguably have been a better design decision. Different observing programs may have different post-observation analysis needs and thus could benefit from multiple post-observation toolsets. On the other hand, any post-observation analysis toolset requires some knowledge of available datasets and the context in which they were obtained, so there’s an advantage to grouping this code with the observing code. It seemed helpful to not splinter the relevant software into too many pieces, so the two components were kept together.

That being said, an important aspect of the design of the post-observation analysis toolkit was a reliance on only the datasets stored in the ATA archive and no additional metadata. The reasoning was that while one might have plenty of expectations about the data on disk from the observing plan or even observing logs, the data on disk are the “ground truth,” and there are always unexpected ways for the logged metadata and recorded visibilities to disagree. There were indeed many cases in practice in which the metadata and the actual datasets disagreed.

For each correlator observing campaign, a table of observing sessions and individual scans is maintained. The session table is populated from the schedule with the program `pdb-stubsess`, with new sessions marked as “unscanned”. After being observed, each session is eventually scanned with `pdb-scansess` and marked as “scanned”. The `pdb-scansess` program creates entries for each scan in the session and in fact reads all of the visibility files completely to search for any potential problems in the data. It also records useful per-scan metadata, the fields of which are listed in Table 2. Of particular note are the `lst0` field, which allows quick analysis of the hour-angle coverage of observations of a source, and the `failtype` field, which records any issues that make the scan unusable. To paraphrase *Anna Karenina*, “Successful observations are all alike; every unsuccessful observation is unsuccessful in its own way.” Thus there’s one value of `failtype`, zero, which indicates success, but a variety of nonzero values indicate possible failure modes, as in the familiar `errno` values returned in Unix error codes.

Name	Type	Units	Description
uid	string		UID of the scan's session
fncode	int		identifies the scan filename within the session data directory via reference to a separate table
ccode	int		identifies the equatorial coordinates of the scan pointing via reference to a separate table
freq	int	MHz	the observing frequency of the scan
focus	int	MHz	the predominant focus setting of the antennas during the scan
tst	int	seconds	the start time of the scan (as a Unix time-since-Epoch)
dur	int	seconds	the duration of the scan
lst0	float	hours	the LST at the beginning of the scan
failtype	int		failure flag if scan is unusable for some reason

Table 2: Metadata recorded in the scan database.

The observation databases are recorded in flat, line-oriented text files using a simple database layer implemented in `pylib/flatdb.py`. While it's probably foolish to implement one's own database layer, the `flatdb` system is simple, fairly efficient, and was not a major sink of programmer time. The motivation for creating it was to take advantage of the fact that the Git repository would effectively provide change tracking and data backup. While a few existing text-based Python database modules were found, they were generally not well-engineered.

Additional utilities were created to populate the observation databases more fully as certain new needs came to light. For instance, `pdb-filldatasizes` computes and records the total data size of each session, to allow reporting of the total survey data volume. `pdb-fillfocus` determines and inserts focus setting information because the importance of this information was not initially obvious.

Several schedule blocks were used to observe AGILITE sources outside of the commensal observing system. Several tools were written to retroactively populate the database with information from these sessions so that all of the relevant information would be centralized. Not all sessions run outside

of the commensal system map onto the schema of the observation database, but many do. The program `misc-scantree` performs a similar task to `pdb-scansess`, printing out processed scan information, but it does not rely on the existence of metadata from `obs-slave.py`. The program `pdb-retrosess` does the job of inserting this information into the database. The two can be linked together in a shell pipeline.

Various other utilities query the databases to analyze, *e.g.*, hour angle coverage of a source, total observing time, or data quality. The most widely-used of these tools is `qdb-dumpsess`, which can summarize the status of an entire observing campaign or one of its component sessions. The simultaneous use of two correlators complicates some queries, since separate scan records are created for each correlator’s data stream. For instance, if one naïvely adds up the integration time of a group of scans on a particular source, the total will be about twice the actual integration time, because two correlators were active. A different database schema could certainly trade off this particular issue for other ones.

A final group of tools integrate the information in the observation database to ARF, the ATA Reduction Framework, the system currently used to analyze the commensal correlator observations. While a discussion of ARF is outside the scope of this document, I’ll mention that the programs `rx-create` and `rx-recreate` stub out ARF “reduction specifications” and link ARF work directories with the appropriate raw datasets.

## 4 Recommendations

I conclude with some recommendations to be kept in mind when designing and implementing future commensal observing campaigns. These are naturally targeted toward projects similar to the one described here and won’t apply to every campaign that could be described as “commensal.”

- **KISS: Keep It Simple, Stupid.** Perhaps the only universal engineering maxim, and it’s as relevant in the commensal context as it is everywhere else. It’s almost always better to get something small working and build out. We certainly had many ideas for the campaign described in this memo that, in retrospect, would have been complete wastes of time to implement.

The following group of recommendations is actually more specific *large* projects than *commensal* projects. For a more authoritative perspective, [Kleinman et al. \(2008\)](#) present some lessons learned from the operations of the Sloan Digital Sky Survey.

- **It's all about the data.** You'll be able to write new software, but you won't be able to take new data, so get the data right. The highest priority, of course, is getting whatever allows you to accomplish your science goals. Beyond that, the more uniform your data are, the easier processing will be — so not only is it important to get the data right, but it makes life a lot easier to think about these things hard before the campaign even starts. Large quantities of data are extremely slow to move around, so production-scale processing needs to require as little copying as possible.
- **Get an end-to-end processing pipeline online as soon as possible.** You don't want to wait until after half your observing is done to realize that you need a new kind of calibration observation, or you're missing some very useful metadata, or something's wrong in the archives. Once a pipeline is in place, you can also start inserting sanity checks to discover bad data just days, not months, after they start coming in. Start out by stubbing out as much as possible (cf. KISS) and fill in the details as you can.
- **Define and monitor observing efficiency metrics.** You want to know if you're on track to reach your science goals, and your observations will almost surely be less efficient in practice than in theory. Choose a just few key metrics to avoid information overload. As with the processing pipeline, the earlier these can be put into regular use, the better.
- **Schedule a time to step back and review.** If observations are running near-continuously, it becomes difficult to take the time to review how the campaign has progressed and ponder how efficiency might be improved. After a large project has gotten rolling, it's probably worthwhile to fall behind on data processing in order to spend a week or so conducting such a review.

These recommendations are more specific to commensal campaigns:



- **“Electronics before concrete.”** This is a slogan promoting cost-effective design adopted by Swiss railroad planners ([Schwager, 2008](#)). The idea is that it’s much cheaper and faster to retrofit existing systems (new signaling systems on existing lines, in the Swiss context) than it is to build new ones from scratch. This is certainly also true when building the infrastructure to run a commensal campaign: you should take advantage of existing infrastructure for non-commensal observations. It only took a few weeks to bolt a small telemetry module onto the existing SETI Prelude system; meanwhile, it took several years to get SonATA, the from-scratch Prelude replacement, on the air.
- **“Organization before electronics before concrete.”** This variation is used by some German planners ([Baumgartner et al., 2011](#)). Their point is large projects often involve multiple unaffiliated actors (rail transit agencies) whose turf wars and internal preferences can lead to plans that are *much* costlier than what would be arrived at by an apolitical efficiency-focused team; thus, when planning a project, one of the most useful things you can do is constructively tackle points of conflict, even though it’s always tempting to avoid them. In a commensal campaign, of course, there are also multiple actors with diverging interests. It was possible to design the telemetry system (§3.3) in such a reliable way only because decisions about observing strategy were negotiated in person and not left to be implicitly made in software.
- **Share as much infrastructure as possible.** This is related to the previous item. In a commensal context, multiple systems doing the same job will find a way to get out of sync and cause problems. The scheduling system of this memo is an example of this. It isn’t bad considered on its own, but its interactions with the SETI and HCRO systems are problematic. Much time has been spent emailing back and forth to negotiate and confirm schedule changes. There’s no good reason for the correlator and beamformer observations to be scheduled separately. The *bad* reason for this is that the two sets of observations are run on separate computer systems and so it is difficult for them to share information. I have no doubt, however, that some arrangement could have been arrived at, avoiding not only the tedious emailing but also making observing more robust thanks to the elimination of *an entire class* of potential failures.

- **Make standalone pilot observations, and expect to perform standalone “patch-up” observations.** It will help the campaign design process if you’ve performed a few test observations without the commensal infrastructure to check out the mechanics and hopefully learn about any potential pitfalls. There will almost definitely be *some* kind of observation that would be good to get that won’t be made (or will be botched) during regular operations, so plan to make occasional standalone observations to cover these gaps. (Don’t forget to propose for the time to make these observations, if necessary!)
- **Be considerate of observatory staff.** This should go without saying. In the particular case of commensal campaigns, hacks to the observatory infrastructure will likely be necessary, and it’s vital that these occur in a fashion acceptable and comprehensible to the staff who will have to deal with them.

Finally, these recommendations may be relevant to the software implementation of a commensal campaign:

- **Avoid bidirectional communications.** There are many more ways to implement a one-way communications channel than a two-way channel. The one-way telemetry stream described in this memo (§3.3) was straightforward to implement and robust, as demonstrated by the smooth replacement of Prelude with SonATA and the easy creation of simple tools such as `obs-master.py` or `obs-listener.py`. The success of the synchronization protocol can be contrasted with that of JSDA, the messaging system used to interlink the ATA subsystems (Gutierrez-Kraybill et al., 2010). JSDA uses a two-way “remote procedure call” model for communications. While it undoubtedly offers important functionality, the JSDA system is complicated, and obscure lockups or communications failures have not been uncommon. The heaviness of the protocol also makes it time-consuming to integrate new software into the messaging network.
- **Use stateless protocols.** Software and hardware will inevitably fail, and testing of your observing systems is likely to be inexhaustive. Using a stateless architecture makes it so you don’t need to even try to handle a variety of tricky problems.

Speaking as the person who drove the commensal correlator system and who has to reduce all of the data it generated, I feel that it performed well, and that this was largely due to some good design decisions early on in the commensal campaign. It's not likely that much of the particular implementation will be portable to future facilities, but I hope that some of these recommendations will be.

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## A Commensal Command Summaries

Name	Description
commensal-mc	Monitor & control of commensal obs by obs user
commensal-mc-helper	Helper for the above
misc-archdir	Print the /ataarchive directory for a session's data
misc-crsync	Copy session data from strato to cosmic
misc-decodeunixtime	Print a Unix time in human-friendly format
misc-diffunixtime	Print the difference between two Unix times in a human-friendly format
misc-dsetnames	Print the names of datasets including a particular source
misc-latestdircmd	Print an eval'able command to change to the directory containing the latest observations
misc-makeunixtime	Convert calendar information into a Unix time
misc-parseplog	Summarize an ongoing commensal observation from its pointing logfile
misc-scantree	Scan an arbitrary tree of visibility data
misc-sessalias	Given a session UID, print its alias
misc-sessuid	Given a session alias, print its UID
misc-whatsnear	Given coordinates, find the nearest observed pointings
obs-hourly	The hourly observing launcher
obs-kickstart	Launch observations right now
obs-launcher	Backend program to launch observations
obs-launcher-helper	Helper for the above
obs-listener.py	Debug the commensal synchronization protocol
obs-master.py	Drive the array simplistically
obs-slave.py	Perform commensal correlator observations
pdb-datealias	Set session aliases from their observing dates
pdb-filldatasizes	Fill in data size information for sessions missing it
pdb-fillfocus	Fill in focus information for sessions missing it
pdb-fillsts	Fill in LST information for sessions missing it

Name	Description
pdb-fillobscfgs	Fill in observing frequency/focus configuration information for sessions missing it
pdb-importagcts	Import database information from the first-generation AGCTS database
pdb-importproj	Import information from one commensal observing project (campaign) into another
pdb-retrosess	Retroactively import the scans and metadata for a session
pdb-scansess	Import scan information for a given session
pdb-stubsess	Stub out session information in the database from the current schedule
plan-archdir	Print out the expected archive directory of a scheduled observing session
plan-importtxt	Import the ATA textual schedule into the commensal schedule
plan-schedsess	Manually schedule an observing block
plan-showsched	Print the current observing schedule
qdb-dumpsess	Print information about an observing campaign or session
qdb-sesssrchacov	Show how well each session covers a given source in a given hour angle range
qdb-srccov	Summarize the coverage of a particular source in a project on a session-by-session basis
qdb-srchacov	Show the hour angle coverage of a given source over the course of the campaign
qdb-summarize	Summarize the overall observing statistics of a campaign
qdb-toscan	Print out a list of sessions that probably need to be scanned
rx-create	Create an ARF reduction specification and workspace for a given session
rx-recreate	Re-realize an ARF reduction workspace for a given session
rx-status	Set the reduction status of a session
rx-suggest	Suggest a useful session to reduce

## B Commensal Python Module Summaries

Name	Description
<code>catalog</code>	Loading and using source catalogs
<code>files</code>	Computing paths for accessing file resources
<code>flatdb</code>	Simple line-oriented textual database
<code>projdb</code>	Session and scan database implementation
<code>protocol</code>	Synchronization protocol implementation
<code>sched</code>	Loading and using observing schedules
<code>util</code>	Miscellaneous utilities